

Fast, Secure Failover for IP

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Outline

- Problem space
- Introduction to ILNP
- IPsec failover solution using ILNP:
 - architecture
 - evaluation
- Summary
- Questions

Problem space

Graceful failover for IPsec

- **IPsec graceful failover currently not defined:**
 - no standard mechanism
 - proprietary solutions exist
 - (IPsec is the basis for military HAIPE IP encryptors)
- Graceful:
 - IPsec session should not be interrupted
 - interruption of IPsec traffic should be minimal
 - failover should be invisible to end-systems

IPsec today

- IPsec Security Association (SA) binds to an IP address:
 - IP address has topological semantics
 - IP address is used as a form of identity
- Failover:
 - connectivity change may involve change of IP address range (new IP routing prefix)
 - change in IP prefix changes end-system state for IPsec (and also for upper layer protocols, e.g. TCP)

End system state for an IP session

- Example: end-system state in TCP/IP
- IP addresses, A
- Transport-layer Port numbers, P
- Changes to interface address binding has impact higher up the stack



Current naming architecture

Protocol Layer	IP
Application	FQDN or IP address
Transport	IP address (+ port number)
Network	IP address
(Interface)	IP address

Entanglement ☹️

FQDN = fully qualified domain name

Introduction to the Identifier Locator Network Protocol (ILNP)



New naming architecture: IP vs ILNP

Protocol Layer	IP	ILNP
Application	FQDN or IP address	FQDN (RFC1958)
Transport	IP address (+ port number)	(Node) Identifier (+ port number)
Network	IP address	Locator
(Interface)	IP address	(dynamic mapping)

Entanglement ☹️

Separation 😊

FQDN = fully qualified domain name



End system state with ILNP

- Example: end-system state in TCP/ILNP:
 - within the context of IPv6
- Locator value, L
- Port numbers, P
- Node Identifier (NID) values, I

$\langle tcp : P_x, P_y, I_x, I_y \rangle \langle ilnp : L_x, L_y \rangle \langle i/f : (L_x) \rangle$

$\langle tcp : P_x, P_y, A_x, A_y \rangle \langle ip : A_x, A_y \rangle \langle i/f : A_x \rangle$



Locator/Identifier Split for ILNP

- **Locator, L:**
 - **Topologically significant.**
 - Names a (sub)network
 - Similar to today's **network routing prefix**
 - Used only for routing and forwarding in the core.
- **(Node) Identifier, NID:**
 - **Is not topologically significant.**
 - Names a logical/virtual/physical node, does **not** name an interface.
- **Upper layer protocols bind only to Identifier.**

ILNP: Engineering

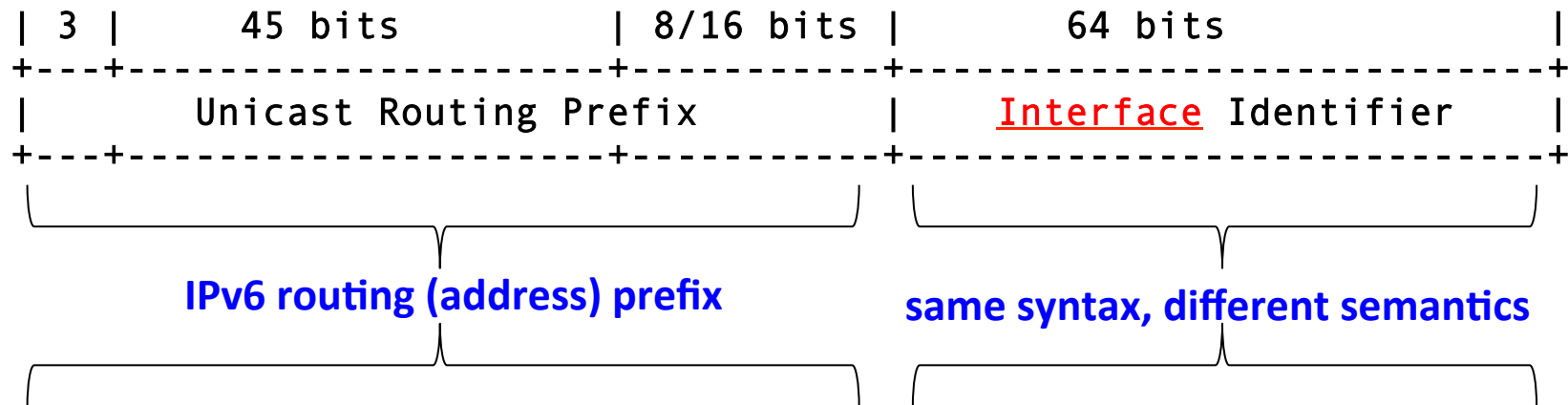
- Main architectural ideas can be applied as extensions to both IPv4 and IPv6:
 - RFCs 6740-6748 (Experimental, IRTF RRG)
- **ILNP extensions to IPv6 → ILNPv6.**
- Non-ILNP nodes see an ordinary IPv6 packet.
- ILNPv6 end-systems see an ILNPv6 packet.
- Focus here is on IPv6, as the engineering is cleaner, but IPv4 is also possible.

ILNPv6

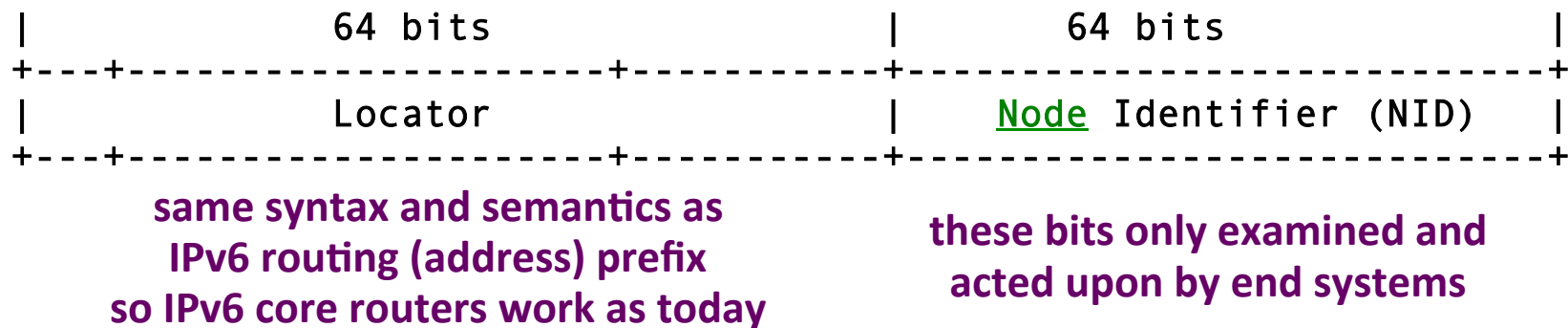
- A set of extensions to IPv6:
 - Same packet format as IPv6, with extensions
 - No changes required in the IPv6 routers
 - Incrementally deployable on IPv6 networks
 - Backwards compatible with IPv6 devices
- Split 128-bit IPv6 address:
 - **64-bit Locator (L64)** -- **(sub)network** name.
 - **64-bit Identifier (NID)** -- **node** name.

IPv6 addresses and ILNPv6

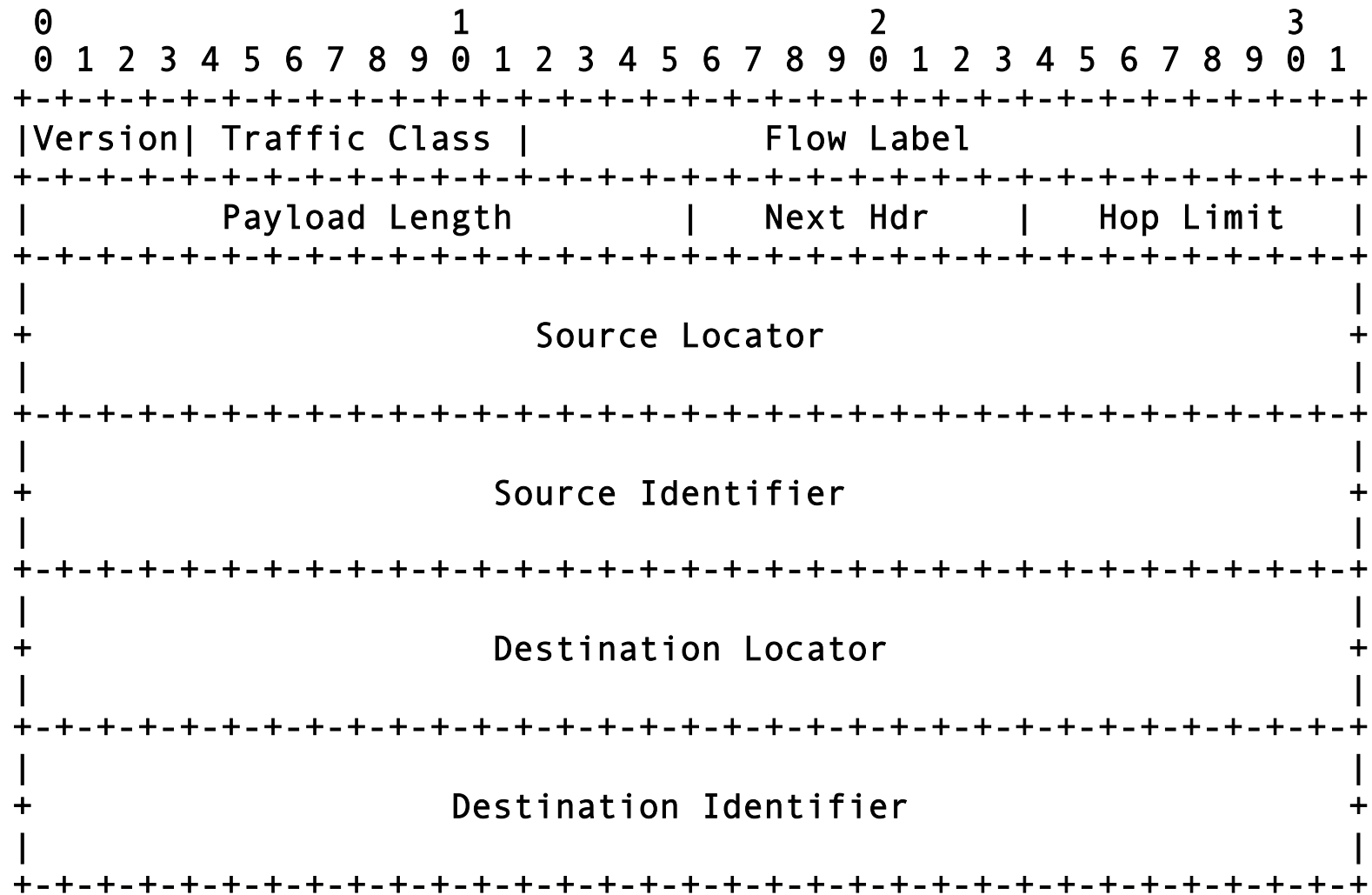
IPv6 (as in RFC3587 + RFC4291):



ILNPv6:



ILNIPv6 packet header (end-system)

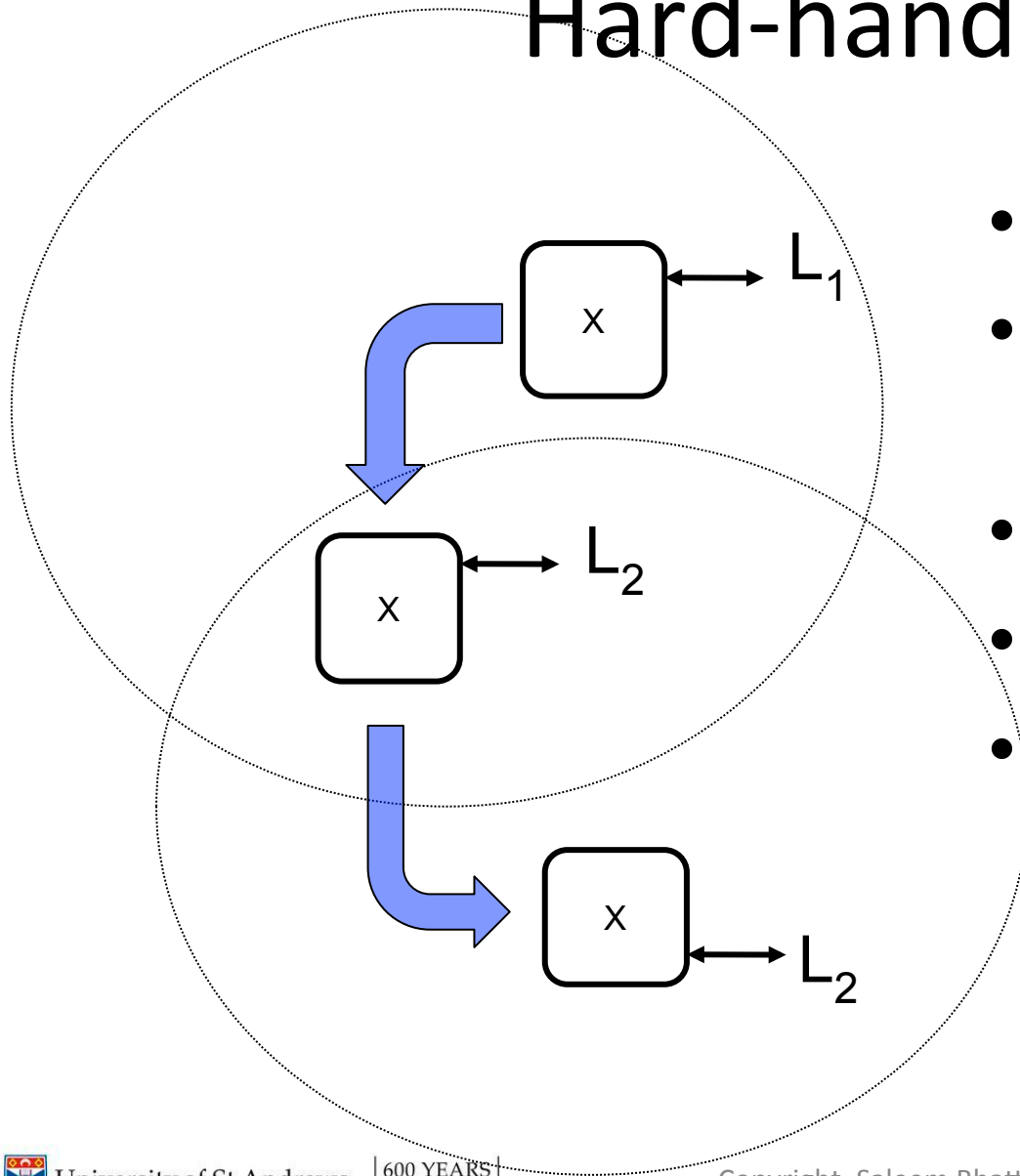


IPsec failover solution using ILNP

IPsec / HAIPE with ILNP Mobility

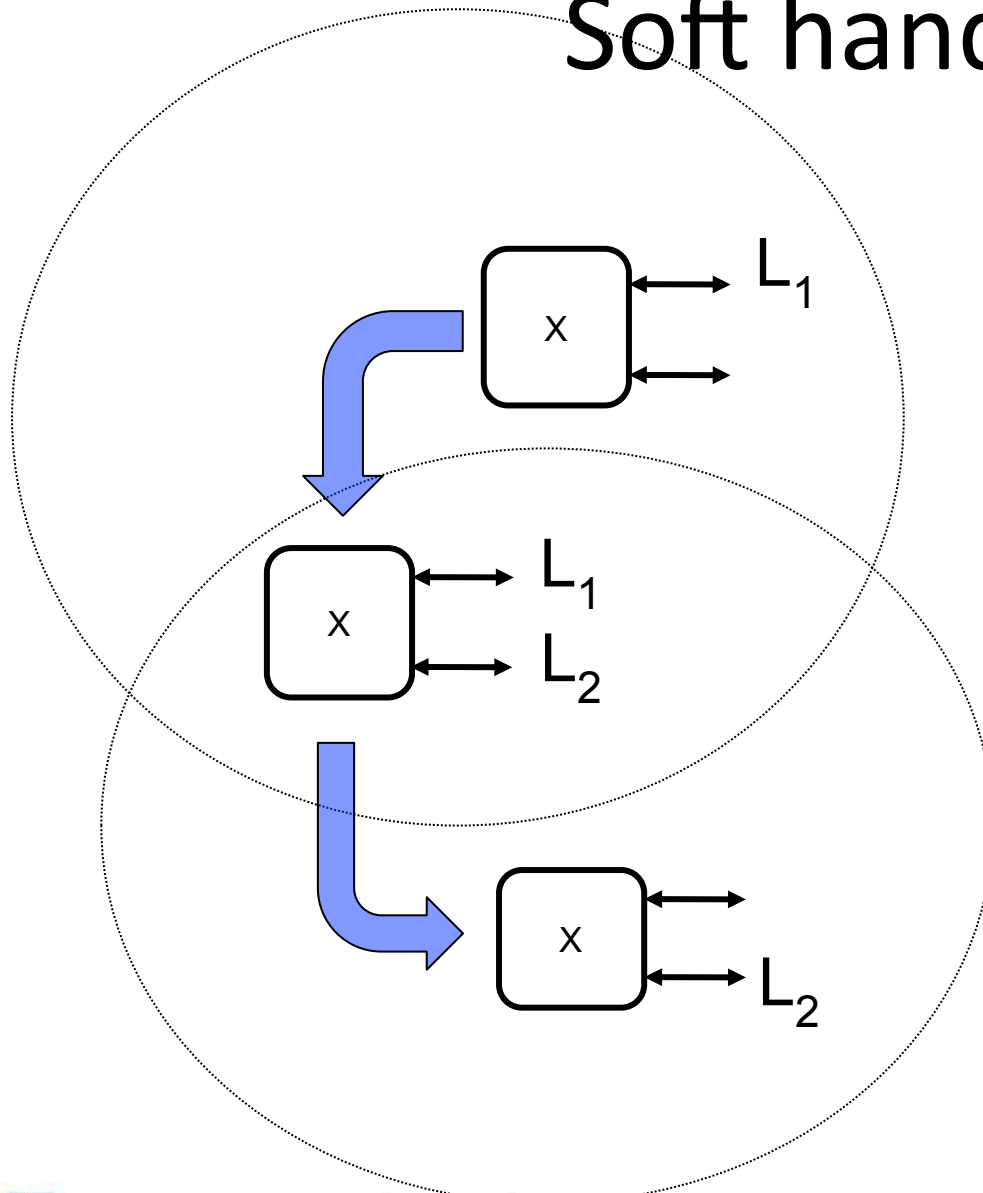
- **IPsec / HAIPE with ILNP binds only to NID:**
 - NID is used in transport layer state
 - NID is not topologically significant
 - preserves end-to-end semantics
- **NID-L binding change does not impact IPsec:**
 - dynamic update of NID-L binding similar to that for Mobile IPv6
 - change in NID-L binding does not impact transport layer sessions
 - Hence, IPsec can be used end-to-end during failover

Hard-handoff [1]



- Start with L_1
- Change IP-level connectivity
- Change L value to L_2
- Use new L_2 value
- Possible packet loss while correspondent node does know about L_2

Soft handoff [1]

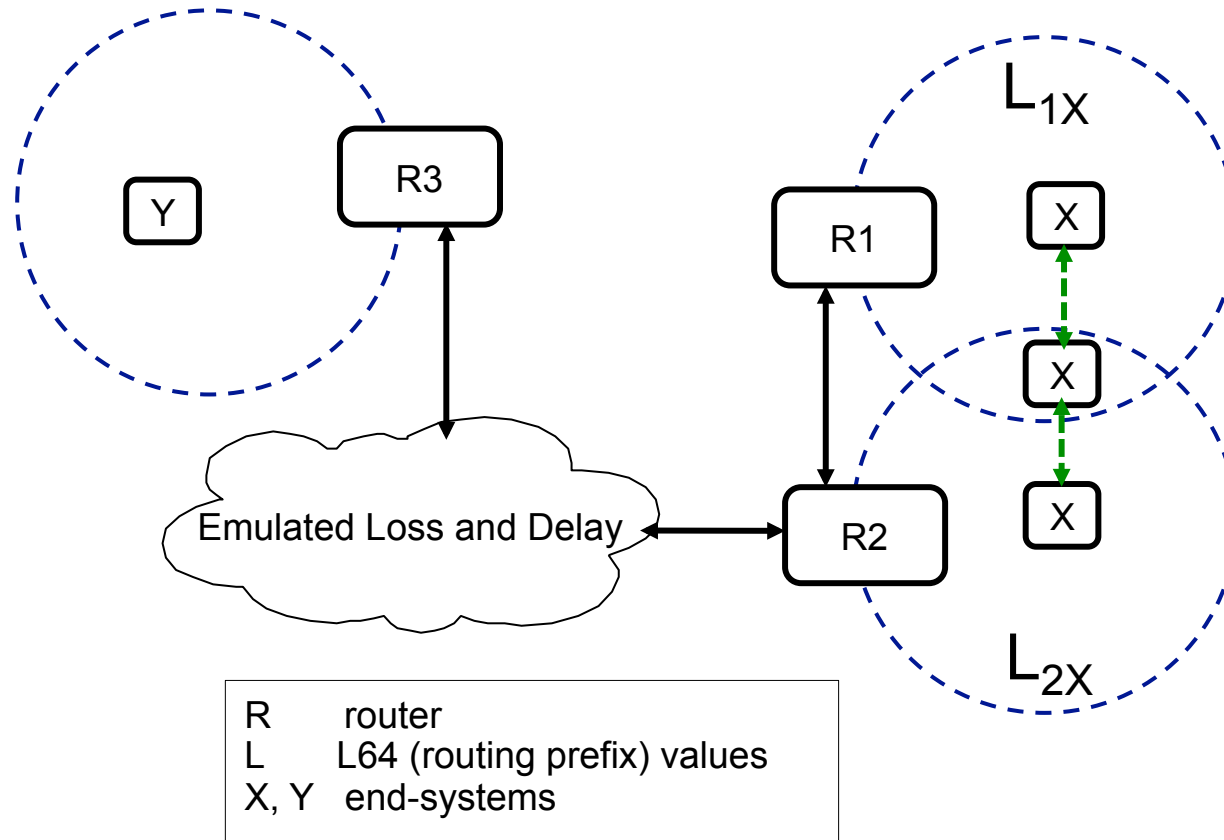


- Start with L_1
- In overlap of cells, use both L_1 and L_2
- Then, handoff to L_2
- Multihoming during handoff
- **Network-layer soft handoff.**

ILNP IPsec advantages

- Completely end-to-end model
- No middleboxes/proxies required – such boxes can be:
 - a single point of failure
 - performance bottleneck
 - point of security attack
- No tunnelling:
 - apart from that which might already be used by IPsec
- No routing changes required:
 - trust boundaries confined to end-systems/sites

Testbed evaluation



Traffic emulation and metrics

TABLE IV
APPLICATION TRAFFIC EMULATION, FLOWS LASTED 65s.

Description	Data Rate [Kbps]	Pkt Size [bytes]	Ref.
Skype / VoIP	64	300	[55], [56]
YouTube / ViIP	658 ^a	1400	[57]

^a This is slightly more than the 632Kbps reported in [57].

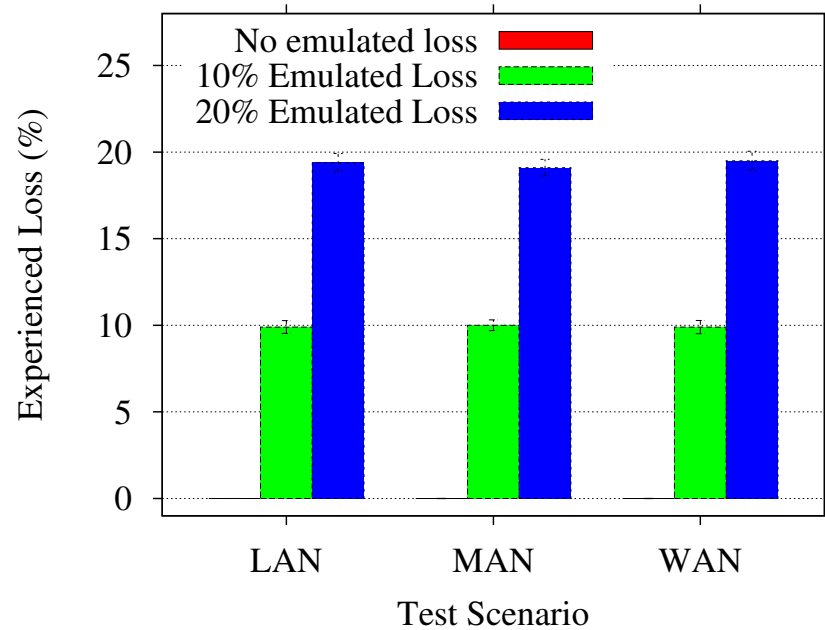
TABLE V
METRICS USED FOR PERFORMANCE EVALUATION.

Metric	Units	Summary	Fig. ^a
loss	%	Application-layer (STP)	4a
f-delay	ms	Time to complete failover	4b
f-overhead	–	# LU/LU-ACK handshakes	4c

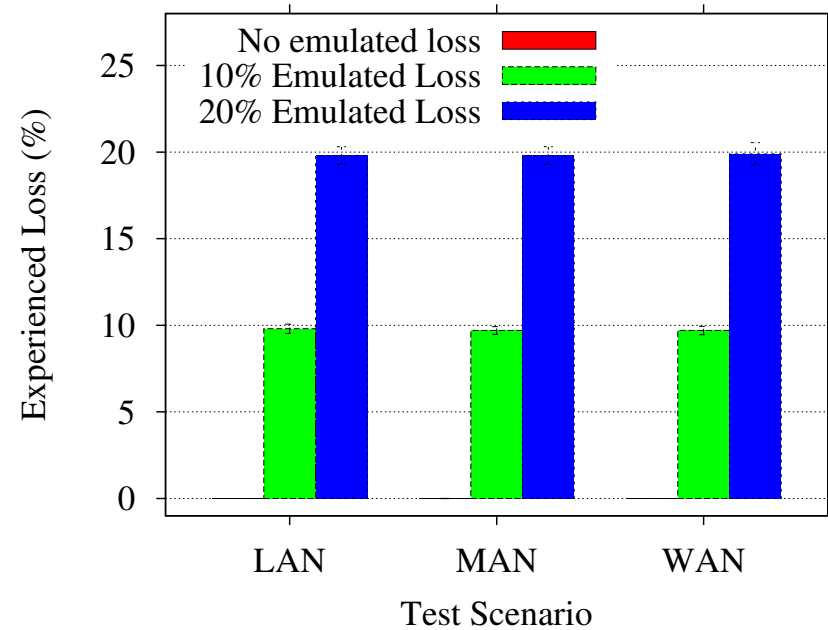
^a The figure showing the related testbed measurements results.

Application layer loss

The mean packet loss of VoIP traffic

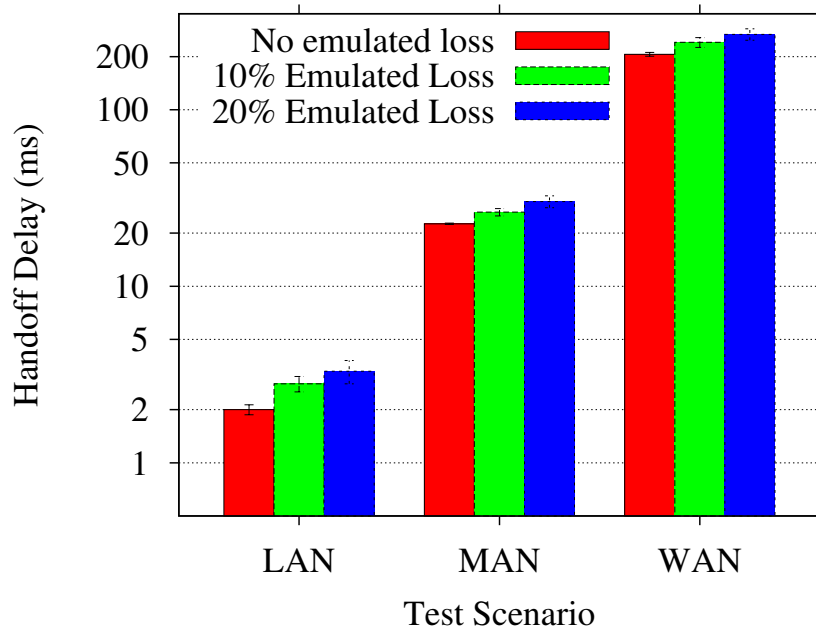


The mean packet loss of Video traffic

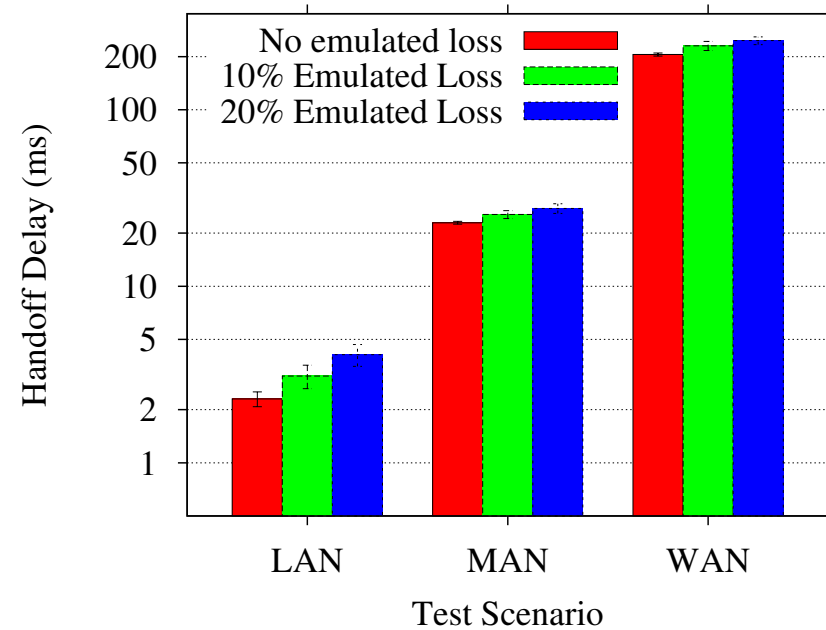


Failover delay

The mean hand-off delay of VoIP traffic

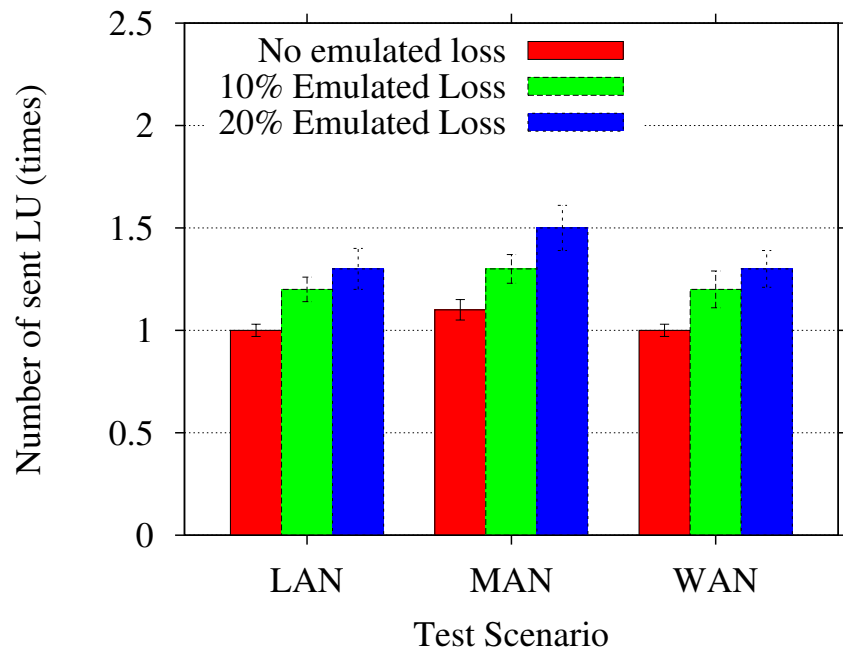


The mean hand-off delay of Video traffic

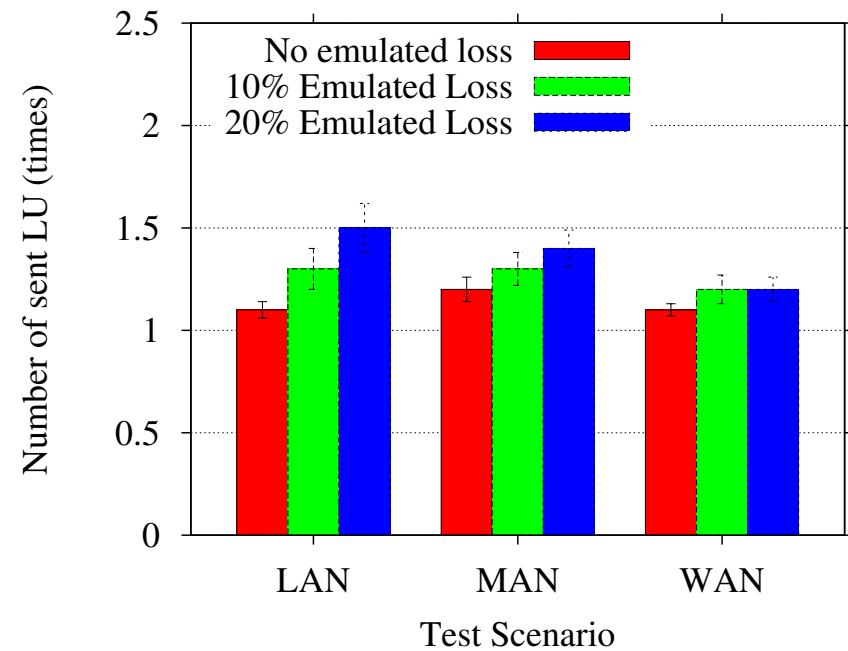


Failover overhead

The mean sent LU per hand-off of VoIP traffic



The mean sent LU per hand-off of Video traffic



Summary

ILNP IPsec failover

- IPsec failover with ILNP:
 - mobility model
- End-to-end model - advantages:
 - no proxy/middlebox
 - trust boundary
 - no additional attack vectors via proxy
- Performance:
 - virtually zero gratuitous loss with soft handoff
 - low overhead

Contacts

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Backup Slides

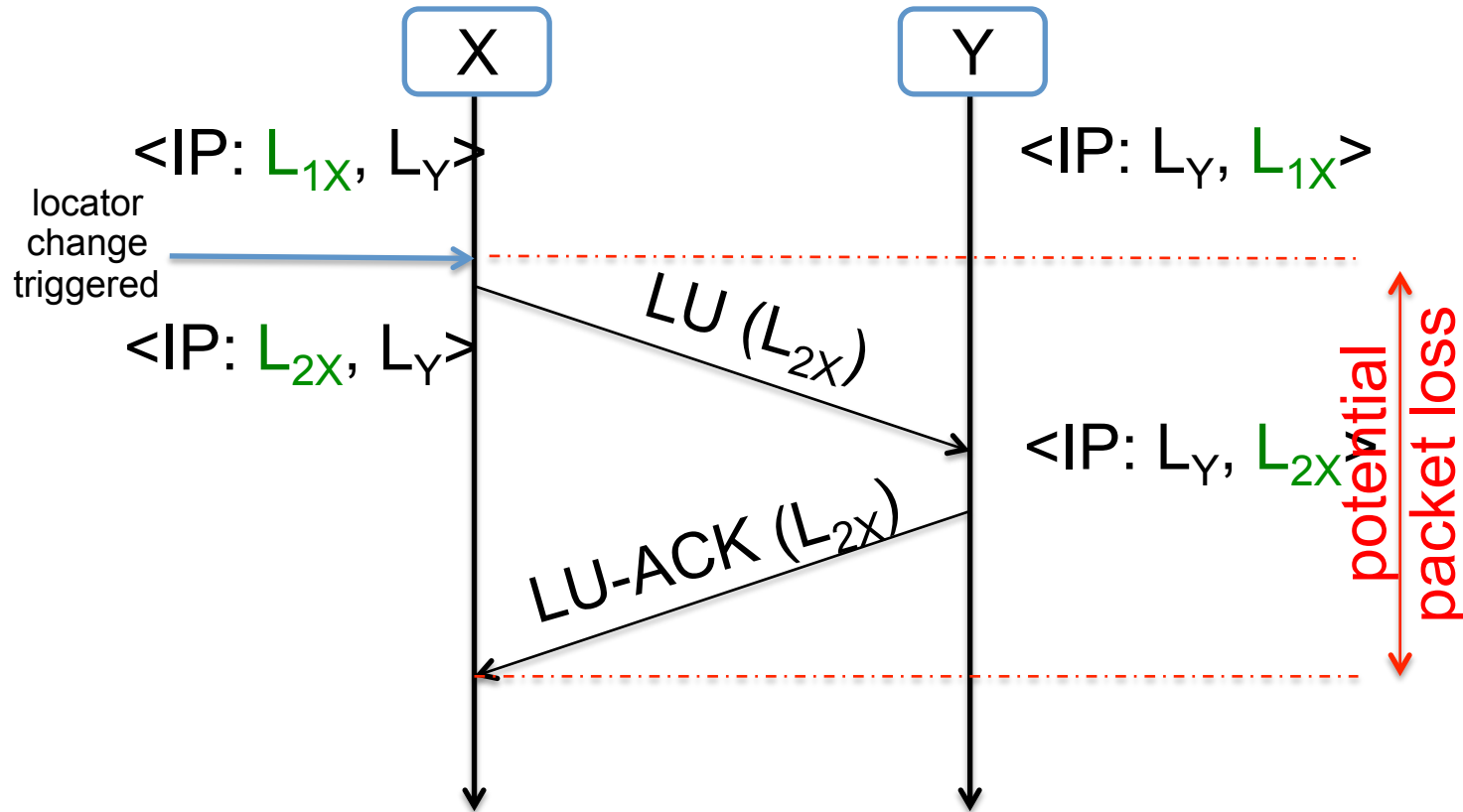
ILNP: Locator Properties

- Locator names an IP Subnetwork.
- Locator is equivalent to an IP Routing Prefix.
- **Nodes can change their Locator values during the lifetime of an ILNP session:**
 - **Enables mobility, multi-homing, NAT, end-to-end IPsec, site-controlled traffic engineering, etc.**
- Multiple Locators can be used simultaneously
 - Enables multi-homing, seamless mobility, end-to-end IPsec, traffic engineering, etc.
- Locators NEVER used by TCP, UDP, SCTP, etc.

ILNP: Identifier Properties

- Identifier names a **node**, not an **interface**
- **Remains constant** during the lifetime of a transport session
 - **Enables IPsec, NAT, & other improvements**
- Nodes have multiple Identifiers concurrently:
 - Only one identifier for a given ILNP session
 - Identifiers are stable over time
- Special NID formats also supported by ILNP:
 - IPv6 Privacy ID extensions, CGAs, etc
- Only NID is used by IPsec, TCP, UDP, SCTP, etc.

Hard-handoff [2]

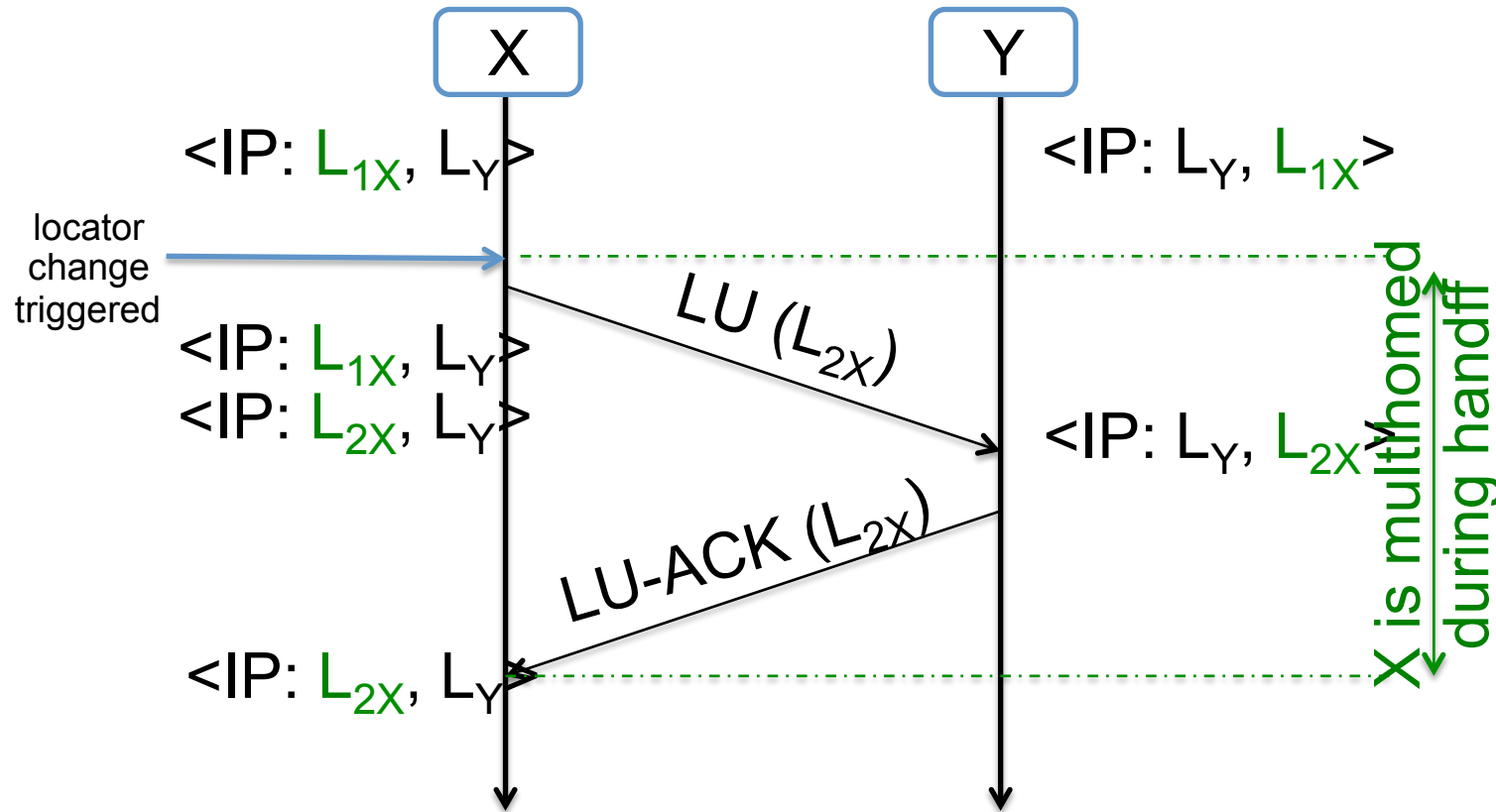


Hard handoff

(similar to Binding Update for Mobile IPv6)

(new L values can be learned from IPv6 router advertisements)

Soft handoff [2]



Soft handoff

(new L values can be learned from IPv6 router advertisements)